Multi-qubit doilies: Enumeration for all ranks and classification for ranks four and five

Axel Muller¹, Metod Saniga², Alain Giorgetti^{*1}, Henri de Boutray³, and Frédéric Holweck^{4,5,6}

¹FEMTO-ST institute, Univ. Bourgogne Franche-Comté, CNRS, France ²Astronomical Institute of the Slovak Academy of Sciences, SK-05960 Tatranská Lomnica, Slovakia ³ColibrITD, France

⁴Université de Technologie de Belfort-Montbéliard, F-90010 Belfort cedex, France ⁵Laboratoire Interdisciplinaire Carnot de Bourgogne (UMR 6303 -CNRS/ICB/UTBM), France

⁶Department of Mathematics and Statistics, Auburn University, Auburn, AL, USA

Abstract

For $N \geq 2$, an N-qubit doily is a doily living in the N-qubit symplectic polar space. These doilies are related to operator-based proofs of quantum contextuality. Following and extending the strategy of Saniga et al. (Mathematics 9 (2021) 2272) that focused exclusively on three-qubit doilies, we first bring forth several formulas giving the number of both linear and quadratic doilies for any N > 2. Then we present an effective algorithm for the generation of all N-qubit doilies. Using this algorithm for N = 4 and N = 5, we provide a classification of N-qubit doilies in terms of types of observables they feature and number of negative lines they are endowed with. We also list several distinguished findings about N-qubit doilies that are absent in the three-qubit case, point out a couple of specific features exhibited by linear doilies and outline some prospective extensions of our approach.

1 Introduction

The doily is a remarkable piece of finite geometry that occurs in a number of disguises. Here, we mention the most prominent ones.

1. The doily as a duad-syntheme geometry. Let us recall a famous Sylvester's construction of the doily [1]. Given a six-element set $M_6 \equiv \{1, 2, 3, 4, 5, 6\}$, a duad is an unordered pair $(ij) \in M_6$, $i \neq j$, and a syntheme is a set of three pairwise disjoint duads, i.e. a set $\{(ij), (kl), (mn)\}$ where $i, j, k, l, m, n \in M_6$ are all distinct. The point-line incidence structure whose points are duads and whose lines are synthemes, with incidence being inclusion, is isomorphic to the doily, as also illustrated in Figure 1.

^{*}Corresponding author, alain.giorgetti@femto-st.fr



Figure 1: A duad-syntheme model of the doily.

- 2. The doily as the Cremona-Richmond configuration. It is a particular 15₃-configuration, i.e. a self-dual configuration of 15 points and 15 lines, with three points on a line and, dually, three lines through a point such that it contains no triangles [2, 3]. Up to isomorphism, there are altogether 245,342 15₃-configurations, of which only the doily enjoys the property of being triangle-free.
- 3. The doily as a generalized quadrangle. A generalized quadrangle GQ(s, t) of order (s, t) is an incidence structure of points and lines (blocks) where every point is on t + 1 lines (t > 0), and every line contains s + 1 points (s > 0) such that if p is a point and L is a line, p not on L, then there is a unique point q on L such that p and q are collinear. The doily is isomorphic to the unique generalized quadrangle with s = t = 2 [4].
- 4. The doily as a symplectic polar space. Given a d-dimensional projective space PG(d, 2) over the two-elements field $\mathbb{F}_2 = \{0, 1\}$ of modulo-2 arithmetic, a polar space \mathcal{P} in this projective space consists of the projective subspaces that are totally isotropic/singular with respect to a given non-singular bilinear form [5, 6]; PG(d, 2) is called the *ambient projective space* of \mathcal{P} . A projective subspace of maximal dimension in \mathcal{P} is called a generator; all generators have the same (projective) dimension r-1. One calls r the rank of the polar space. The symplectic polar space $\mathcal{W}(2N-1,2), N \geq 1$, consists of all the points of PG(2N-1,2), $\{(x_1, x_2, \ldots, x_{2N}) : x_j \in \{0, 1\}, j \in \{1, 2, \ldots, 2N\}\} \setminus \{(0, 0, \ldots, 0)\}$, together with the totally isotropic subspaces with respect to the standard symplectic form

$$\sigma(x,y) = x_1 y_{N+1} - x_{N+1} y_1 + x_2 y_{N+2} - x_{N+2} y_2 + \dots + x_N y_{2N} - x_{2N} y_N.$$
(1)

Throughout the paper, the space name $\mathcal{W}(2N-1,2)$ is often shortened as W_N . This space features

$$|W_N|_p = 4^N - 1$$

points and

$$|W_N|_q = (2+1)(2^2+1)\cdots(2^N+1)$$

generators. The doily is isomorphic to the symplectic polar space of rank $N = 2, \mathcal{W}(3, 2)$.

5. *Multi-qubit doilies*. This paper is about doilies related to Kochen–Specker operator-based proofs of quantum contextuality, to be called *N*-qubit doilies or multi-qubit doilies. We

follow the terminology and notation of Section 2 of [7], to which the reader can refer for more finite-geometric background. Let

$$X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \quad Y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix} \text{ and } Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

be the Pauli matrices, I the identity matrix, ' \otimes ' the tensor product of matrices and $\mathcal{I}_N \equiv I_{(1)} \otimes I_{(2)} \otimes \ldots \otimes I_{(N)}$, and let $\mathcal{S}_N = \{G_1 \otimes G_2 \otimes \cdots \otimes G_N : G_j \in \{I, X, Y, Z\}, j \in \{1, 2, \ldots, N\}\} \setminus \{\mathcal{I}_N\}$. The $4^N - 1$ N-qubit observables of \mathcal{S}_N can be bijectively identified with the $4^N - 1$ points of $\mathcal{W}(2N - 1, 2)$ in such a way that any two commuting observables are represented by collinear points and the product of the three observables lying on a line of $\mathcal{W}(2N - 1, 2)$ is $+\mathcal{I}_N$ or $-\mathcal{I}_N$ (see, for example, [8, Section 5.3.2]). If the symplectic form in the ambient space PG(2N - 1, 2), defining $\mathcal{W}(2N - 1, 2)$, is given by Eq. (1), then the corresponding bijection reads

$$G_j \leftrightarrow (x_j, x_{j+N}), \ j \in \{1, 2, \dots, N\},\tag{2}$$

with the assumption that

$$I \leftrightarrow (0,0), X \leftrightarrow (0,1), Y \leftrightarrow (1,1), \text{ and } Z \leftrightarrow (1,0).$$
 (3)

To briefly illustrate this property, let us consider the three-qubit $\mathcal{W}(5,2)$ and one of its lines, say (0,1,1;1,1,0), (1,0,0;0,0,1) and (1,1,1;1,1,1). Using the correspondences (2) and (3) we find that the corresponding observables are $X \otimes Y \otimes Z$, $Z \otimes I \otimes X$ and $Y \otimes Y \otimes Y$, respectively; these observables indeed pairwise commute and their product is $+I \otimes I \otimes I$.

In what follows, W_N will always be understood as having its points labeled by the N-qubit observables as described above, and any doily lying in it, together with the inherited labeling, will be called an N-qubit doily ($N \ge 2$). Slightly rephrased, an N-qubit doily is a doily whose points are bijectively identified with 15 specific observables from S_N , such that any two commuting observables share the same line, and, given any line, the product of (any) two observables lying on it is, up to a sign, equal to the remaining observable on it. A line of an N-qubit doily will be called *positive* (resp. *negative*) if the product of its three observables is $+\mathcal{I}_N$ (resp. $-\mathcal{I}_N$). To avoid any possible misunderstanding, it is worth mentioning that the product of observables is the (ordinary) matrix product, denoted by a dot (.), induced by the following multiplication table of Pauli matrices.

From here on, the geometrical points are considered to be finite words on the four-letter alphabet $\{I, X, Y, Z\}$ that encode the observables $G_1 \otimes G_2 \otimes \cdots \otimes G_N$, while omitting the symbol \otimes for the tensor product and forgetting in the sequel about the matrix nature of I, X, Y and Z.

Contributions and paper outline. Our contributions start in Section 2, with a presentation of several facts about N-qubit doilies that motivates the design of an effective algorithm to enumerate N-qubit doilies for any rank N (Section 4). By geometric considerations, we first establish in Section 3 closed formulas for the numbers of N-qubit doilies. As these numbers increase rapidly with N, the enumeration algorithm can in practice only be executed for small numbers of qubits. We use it in Section 5 to classify N-qubit doilies for N = 4 and N = 5, according to their types of observables and their configurations of negative lines. We thus produce precise tables for the number of doilies in each category/class, reproduced in the appendices of this paper. Section 5 also analyzes these results and points out various findings about N-qubit doilies that are absent in the known three-qubit case. Section 6 concludes and outlines some prospective extensions of our approach.

2 Some basic facts about multi-qubit doilies

2.1 Patterns formed by negative lines

It is a straightforward task to work out possible types of patterns of negative lines an N-qubit doily can be endowed with. This classification follows readily from the facts that each grid in the doily must contain an odd number of negative lines and that two different grids have two intersecting lines in common. And as a grid has an even number of lines the types of configurations come in complementary pairs, as depicted in Figure 2.



Figure 2: Generic representatives of the twelve different types of configurations of negative lines (bold) that can be found in a multi-qubit doily.

Let us give a brief description of the individual types of configurations. In Type 3 the three negative lines are pairwise disjoint and lie in a grid; that is, their dual is a tricentric triad. Type 4 features three pairwise disjoint lines not belonging to a grid and their unique transversal. In Type 5 the five negative lines form a pentagon. Type 6 contains the three lines from Type 3 plus three

concurrent lines, whose point of concurrence is not lying on any of the three former lines. Type 7A contains six lines forming a hexagon and a unique line disjoint from any of the six. Type 7B is a particular union of two Types 4 and an extra line or, equivalently, is composed of the five lines of a grid and two disjoint lines. A two-qubit doily features just a Type 3 pattern, while in a three-qubit doily we can find all the patterns from Type 3 to Type 7A inclusive [7].

2.2 Linear and quadratic doilies

Following [7], we will also distinguish between two kinds of doilies, referred to as linear and quadratic. A *linear* N-qubit doily spans a PG(3, 2) of the ambient PG(2N - 1, 2). This means that the three lines of a perp-set of such a doily are coplanar, i. e. lie in a PG(2, 2) of the PG(3, 2), a tricentric triad corresponds to a line of the PG(3, 2) and the plane defined by a unicentric triad of the doily passes through its center. Figure 3 serves as a graphical illustration of these features for N = 4. In the doily we selected a perp-set (blue) and colored the remaining lines red. The model of PG(3, 2) is based on a 3-D tetrahedral model of Polster [9]; our version features all the points but not all the lines of the model in order to avoid too crowded appearance of the figure. The two red points at the side lie on the line passing via IYZI that would be perpendicular to the plane of the drawings. Each black line of the PG(3, 2) is non-isotropic and corresponds to a tricentric triad in the doily.



Figure 3: A linear four-qubit doily with one of its perp-sets highlighted in blue color (a) and the corresponding PG(3,2) of PG(7,2) it spans (b). One can readily see that the three lines of the perp-set lie in a plane of the PG(3,2) and the three points on a non-isotropic line of the space (black) correspond to a tricentric triad of the doily.

A quadratic N-qubit doily spans a PG(4,2) of the ambient PG(2N-1,2), being, in fact, isomorphic to the geometry formed by 15 points and 15 lines lying on a parabolic quadric Q(4,2) in



Figure 4: An illustration of the fact that a perp-set of a quadratic (four-qubit) doily (a) spans a PG(3,2) (b) of the PG(4,2) spanned by the doily.

this PG(4,2). This quadric, as any other parabolic quadric in PG(4,2), has a remarkable property that all its tangent hyperplanes pass through the same point J, called the nucleus (see, e.g., [5]). Any tricentric triad of such a doily defines a plane in the PG(4,2) that contains J; a unicentric triad also defines a plane, this plane passing through the remaining third point lying on the line defined by J and the (unique) center of the triad. Moreover, all the 15 PG(3,2)s passing through J intersect our quadric in three concurrent lines that form a perp-set of the doily. Figure 4 offers a pictorial illustration of some of these properties. We again take a four-qubit doily, where we highlighted a perp-set (blue). Now the three lines of the perp-set are not coplanar as in the case of linear doily, but span a PG(3,2). We colored the remaining eight points (and the totally-isotropic lines) of the PG(3,2) in yellow in order to stress the property that the only points shared by the doily and this PG(3,2) are the (blue) points of the perp-set. There are two "distinguished" points of the PG(3,2), namely ZYXX and ZYIX, which lie on the remaining seventh line passing via IIXI; the point ZYXX is nothing but the nucleus of the parabolic quadric our particular doily is located on. Given a perp-set, we know that there are four tricentric and four unicentric triads contained in it. In our particular perp, the four tricentric triads are {XXIX, XZII, ZIXI}, {XXIX, XZXI, ZIII}, $\{XZXI, XXXX, ZIXI\}$ and $\{XXXX, ZIII, XZII\}$; one can readily check that the product of the three observables in any of them is ZYXX (the nucleus). The four *unicentric* triads of our perp-set are {XZXI, XXXX, ZIII}, {XZXI, ZIXI, XXIX}, {ZIII, XZII, XXIX} and $\{XXX, XZII, ZIXI\}$; the product of the observables in any of them is ZYIX, i.e. the second distinguished point. By this construction we get a (different) PG(3,2) for any of the 15 perp-sets of the doily; and because in any of these perp-sets the four tricentric triads always define the nucleus, ZYXX, we get altogether 15 PG(3,2)s that share the point ZYXX, these 15 spaces lying in that particular PG(4,2) of the ambient PG(7,2) that contains the quadric of our selected doily.

In a recent paper [7], four of the authors have thoroughly analyzed and classified three-qubit doilies. To this end, they first explicitly computed all 63 perp-sets, 36 hyperbolic quadrics and 28 elliptic quadrics living in $\mathcal{W}(5,2)$. Then, employing the fact that a linear doily is isomorphic to the intersection of two perp-sets with non-collinear nuclei, they computed and classified all $63 \times 32/3! = 336$ linear doilies of the $\mathcal{W}(5,2)$. In the next step, making use of the property that a quadratic doily is isomorphic to the intersection of an elliptic quadric and a hyperbolic quadric, they generated and classified all $36 \times 28 = 1008$ quadratic doilies of the $\mathcal{W}(5,2)$. The procedure described above is, however, not a viable one for N > 3, as we would first need to compute all $\mathcal{W}(5,2)$ s living in a particular $\mathcal{W}(2N-1,2)$, N > 3, and then in each of them compute 336 linear and 1008 quadratic doilies following the strategy of [7]. Instead, we shall follow (in Section 4) a different, and reasonably faster, approach that makes use of some properties of an ovoid of a doily. In particular, we shall start with a particular N-qubit ovoid, i.e. a set of five mutually anticommuting N-qubit observables whose product is $\pm \mathcal{I}_N$, and introduce a unique algebro-geometrical recipe with the help of which one can find all the N-qubit doilies having this particular ovoid in common. Before embarking on this path, however, we shall introduce several general formulas for the number of both linear and quadratic doilies of $\mathcal{W}(2N-1,2)$, valid for any N > 2, so that we already have certain important numbers at hand to validate some of our subsequent, mostly computer-assisted, results.

2.3 Contextuality degree

All multi-qubit doilies are observable-based proofs of the Kochen–Specker theorem, that establishes that no Non-Contextual Hidden Variables (NCHV) model can reproduce the outcomes of quantum mechanics. This contextuality property is related to a linear problem, as follows. Let A be the incidence matrix of the points on the lines of a finite geometry, such as the doily. Its coefficients are in the two-elements field $\mathbb{F}_2 = \{0, 1\}$, its l rows correspond to the geometric lines and its pcolumns to the geometric points (for the doily, l = p = 15). The positive (resp. negative) nature of a line is encoded by a 0 (resp. 1) for the corresponding coefficient of the valuation vector Ein \mathbb{F}_2^l . Then a quantum geometry is contextual iff there is no vector x such that Ax = E. The contextuality degree is the minimal Hamming distance between a vector Ax and the vector E [10]. The contextuality degree is the minimal number of line valuations that one should change to make the quantum geometry satisfiable by an NCHV model.

Proposition 1. All multi-qubit doilies have a contextuality degree of 3.

Proof. All multi-qubit doilies have the same incidence matrix A. Accordingly, the only parameter that is changing between all the doilies is the vector E, which only depends on the configuration of their negative lines. We have seen that there are only 12 such configurations. For each of these 12 configurations, we have computed the Hamming distance between Ax and E, for all vectors x in \mathbb{F}_2^{15} . It turns out that the minimal Hamming distance is always 3.

In practice, we did not write by hand the 12 possible E vectors, but we computed these vectors from the 5-qubit doilies, because, as described later, we have checked by enumeration that these doilies present all the configurations.

3 Numbers of multi-qubit doilies

This section proposes and justifies closed formulas for the numbers of linear and quadratic doilies in $\mathcal{W}(2N-1,2)$. Before all we introduce some well-known formulas. First, we introduce the Gaussian (binomial) coefficient

$$\begin{bmatrix} n \\ k \end{bmatrix}_{q} = \prod_{i=1}^{k} \frac{q^{n-k+i}-1}{q^{i}-1} = \frac{(q^{n}-1)\dots(q^{n-k+1}-1)}{(q^{k}-1)\dots(q-1)}$$
(4)

where $0 \le k \le n$ and q is a power of a prime, which gives the number of subspaces of dimension¹ k-1 in a projective space PG(n-1,q) of dimension n-1 over \mathbb{F}_q . More generally, the number of (k-1)-dimensional spaces of PG(n-1,q) that pass through a fixed (l-1)-dimensional space is

$$\begin{bmatrix} n-l\\k-l \end{bmatrix}_q.$$
(5)

Next, for a symplectic polar space $\mathcal{W}(2N-1,q)$ embedded in a projective space $\mathrm{PG}(2N-1,q)$, the number of its k-dimensional spaces is given by (see, e.g., [11, Lemma 2.10])

$$\begin{bmatrix} N\\k+1 \end{bmatrix}_{q} \prod_{i=1}^{k+1} \left(q^{N+1-i} + 1 \right)$$
(6)

and the number of k-dimensional spaces through a fixed m-dimensional space [11, Corollary 2.11] equals

$$\binom{N-m-1}{k-m}_{q} \prod_{i=1}^{k-m} \left(q^{N-m-i} + 1 \right).$$
 (7)

Further, let \perp be a symplectic polarity of $\operatorname{PG}(n,q)$ and let denote by S^{\perp} the polar space of a subspace S. If S is of dimension k, then S^{\perp} has dimension n - k - 1. A projective subspace S of $\operatorname{PG}(n,q)$ is called *isotropic* if $S \cap S^{\perp} \neq \emptyset$ and *non-isotropic* if $S \cap S^{\perp} = \emptyset$. An isotropic S is called *totally isotropic* if $S \subseteq S^{\perp}$. It is easy to see that if S is a totally isotropic subspace, then every subspace contained in S is also totally isotropic. Moreover,

$$S \subseteq T^{\perp} \Rightarrow T \subseteq S^{\perp}.$$
(8)

In order to prove the two theorems below, we will need a couple of lemmas.

Lemma 2. If a PG(3,2) of the ambient PG(5,2) equipped with a symplectic polarity \perp contains a totally-isotropic PG(2,2), then it contains exactly three such PG(2,2)s, passing through a common (totally-isotropic) PG(1,2).

Proof. First, there are no totally-isotropic PG(3, 2)s in the PG(5, 2). Given a totally-isotropic PG(1, 2) of PG(5, 2), S, there are (see Eq. (7) for q = 2, N = 3, k = 2 and m = 1) three totally-isotropic PG(2, 2)s passing through it. Denoting these as T_i^{\perp} (i = 1, 2, 3), the lemma then follows from the fact that $S^{\perp} \cong PG(3, 2)$, $T_i^{\perp} = T_i$, and property (8).

¹All dimensions in this section are projective dimensions.

Remark 3. For N > 3, PG(2N - 1, 2) features also totally-isotropic PG(3, 2)s; any other of its PG(3, 2)s endowed with totally-isotropic PG(2, 2)s has the property as described in Lemma 2.

Lemma 4. If a PG(4,2) of the ambient PG(7,2) equipped with a symplectic polarity \perp contains a totally-isotropic PG(3,2), then it contains exactly three such PG(3,2)s, passing through a common (totally-isotropic) PG(2,2).

Proof. The proof parallels that of the preceding lemma. First, there are no totally-isotropic PG(4,2)s in the PG(7,2). Given a totally-isotropic PG(2,2) of PG(7,2), S, there are (see Eq. (7) for q = 2, N = 4, k = 3 and m = 2) three totally-isotropic PG(3,2)s passing through it. Denoting these as T_i^{\perp} (i = 1, 2, 3), the lemma then follows from the fact that $S^{\perp} \cong PG(4,2), T_i^{\perp} = T_i$, and property (8).

Remark 5. For N > 4, PG(2N - 1, 2) features also totally-isotropic PG(4, 2)s; any other of its PG(4, 2)s endowed with totally-isotropic PG(3, 2)s has the property as described in Lemma 4.

Next, through a (totally-isotropic) point of PG(5, 2), S, there pass 15 totally-isotropic PG(1, 2)s, T_j^{\perp} (j = 1, 2, 3, ..., 15) and the same number of PG(2, 2)s. Given the facts that $S^{\perp} \cong PG(4, 2)$ and $T_j \cong PG(3, 2)$, a PG(4, 2) of PG(5, 2) will contain 15 PG(3, 2)s of type defined by Lemma 2 concurring at a point, namely the pole of this particular PG(4, 2). As PG(4, 2) contains altogether 31 PG(3, 2)s, each of the remaining 16 PG(3, 2)s does not contain totally-isotropic PG(2, 2)s and so hosts a unique linear doily. As each such doily can be viewed as a projection of a quadratic doily from the pole, a PG(4, 2) is found to be spanned by 16 quadratic doilies.

Remark 6. If a PG(4,2) of the ambient PG(2N - 1, 2), N > 3, is devoid of totally-isotropic PG(3,2)s, then it is of the type described above, i.e. it entails 16 quadratic doilies.

3.1 Number of linear doilies

Theorem 7. For any $N \ge 2$ the number of linear doilies in W(2N-1,2) is

$$D_l(N) = \begin{bmatrix} 2N\\4 \end{bmatrix}_2 - \begin{bmatrix} N\\4 \end{bmatrix}_2 \prod_{i=1}^4 \left(2^{N+1-i}+1\right) - 7\begin{bmatrix} N\\3 \end{bmatrix}_2 2^{2N-6} \prod_{i=1}^3 \left(2^{N+1-i}+1\right)/3.$$
(9)

Proof. A linear doily of $\mathcal{W}(2N-1,2)$ spans a particular PG(3,2) of the ambient PG(2N-1,2) that does not contain any totally-isotropic PG(2,2). And since any such PG(3,2) is spanned by a single linear doily, the number of linear doilies of $\mathcal{W}(2N-1,2)$ is thus equal to the number of PG(3,2)s that are devoid of totally-isotropic planes. To find this number, from Eq. (4) we first note that there are altogether

$$\begin{bmatrix} 2N\\4 \end{bmatrix}_2 \tag{10}$$

PG(3,2)s in PG(2N-1,2), out of which

$$\begin{bmatrix} N\\4 \end{bmatrix}_2 \prod_{i=1}^4 \left(2^{N+1-i} + 1 \right)$$
(11)

(Eq. (6) with k = 3 and q = 2) are totally isotropic.

To ascertain the cardinality of the remaining PG(3, 2)s that feature totally-isotropic PG(2, 2)s, we proceed as follows. We first observe that by Eq. (6) with k = 2 and q = 2 there are

$$\begin{bmatrix} N\\3 \end{bmatrix}_2 \prod_{i=1}^3 \left(2^{N+1-i} + 1 \right)$$
(12)

totally-isotropic PG(2,2)s in PG(2N-1,2). Next, with k = 3 and m = 2 in (7), it follows that there are

$$\begin{bmatrix} N-3\\1 \end{bmatrix}_2 (2^{N-3}+1) = 2^{2(N-3)} - 1$$
(13)

totally-isotropic PG(3, 2)s passing through a totally-isotropic PG(2, 2). And since the total number of PG(3, 2)s passing via a PG(2, 2) of PG(2N - 1, 2) is

$$\begin{bmatrix} 2N-3\\4-3 \end{bmatrix}_2 = 2^{2N-3} - 1 \tag{14}$$

(as stemming from Eq. (5) for n = 2N, k = 4, l = 3 and q = 2), through a totally-isotropic PG(2, 2) there pass

$$2^{2N-3} - 1 - \left(2^{2(N-3)} - 1\right) = 7 \times 2^{2N-6}$$
(15)

isotropic PG(3, 2)s apart from those that are totally isotropic. Hence, the number of those PG(3, 2)s of PG(2N - 1, 2) that are endowed with totally-isotropic PG(2, 2)s – with the exclusion of totally isotropic ones – amounts to $(12) \times (15)/3$, where we also took into account (see Remark 3) that any such PG(3, 2) features just three totally-isotropic PG(2, 2)s. All in all, there are

$$\begin{bmatrix} 2N\\4 \end{bmatrix}_2 - \begin{bmatrix} N\\4 \end{bmatrix}_2 \prod_{i=1}^4 \left(2^{N+1-i}+1\right) - 7\begin{bmatrix} N\\3 \end{bmatrix}_2 2^{2N-6} \prod_{i=1}^3 \left(2^{N+1-i}+1\right)/3$$

PG(3,2)s in the ambient PG(2N-1,2) that are devoid of totally-isotropic PG(2,2)s, and so the same number of linear doilies in $\mathcal{W}(2N-1,2)$.

Given the fact that the three lines of a perp-set of a linear doily span a PG(2, 2), and namely that PG(2, 2) that features just three totally-isotropic PG(1, 2)s, we arrive at the interesting expression

$$D_l(N) = \frac{4}{15} 4^{N-3} \Theta_2(N), \tag{16}$$

for the number of linear doilies in $\mathcal{W}(2N-1,2)$, where

$$\Theta_2(N) = \frac{1}{16} 2^{2N} \prod_{i=1}^2 \frac{2^{N-2+i}-1}{2^i-1} \prod_{i=1}^2 (2^{N+1-i}+1)$$
(17)

is the number of those PG(2, 2)s of the ambient PG(2N - 1, 2) each of which features just three totally-isotropic PG(1, 2)s.

3.2 Number of quadratic doilies

Theorem 8. For any $N \ge 3$ the number of quadratic doilies in W(2N - 1, 2) is

$$D_q(N) = 16\left(\begin{bmatrix} 2N\\5 \end{bmatrix}_2 - \begin{bmatrix} N\\5 \end{bmatrix}_2 \prod_{i=1}^5 \left(2^{N+1-i} + 1 \right) - 15 \begin{bmatrix} N\\4 \end{bmatrix}_2 2^{2N-8} \prod_{i=1}^4 \left(2^{N+1-i} + 1 \right) / 3 \right).$$
(18)

Proof. A quadratic doily of W(2N-1,2) spans a particular PG(4,2) of the ambient PG(2N-1,2) that does not contain any totally-isotropic PG(3,2). And since any such PG(4,2) is spanned by (see Remark 6) 16 such doilies that are all unique to this space, the number of quadratic doilies of W(2N-1,2) is thus equal to 16 times the number of PG(4,2)s that are devoid of totally-isotropic PG(3,2)s. To find the latter number, we again start with Eq. (4) that tells us that there are altogether

$$\begin{bmatrix} 2N\\5 \end{bmatrix}_2 \tag{19}$$

PG(4,2)s in PG(2N-1,2), out of which

$$\begin{bmatrix} N\\5 \end{bmatrix}_2 \prod_{i=1}^5 \left(2^{N+1-i}+1\right) \tag{20}$$

(Eq. (6) with k = 4 and q = 2) are totally isotropic.

To ascertain the cardinality of the remaining isotropic PG(4, 2)s, we proceed as follows. We first observe that by Eq. (6) with k = 3 and q = 2 there are

$$\begin{bmatrix} N\\4 \end{bmatrix}_2 \prod_{i=1}^4 \left(2^{N+1-i} + 1 \right)$$
(21)

totally-isotropic PG(3,2)s in PG(2N-1,2). Next, with k = 4, m = 3 and q = 2 in (7) it follows that there are

$$\begin{bmatrix} N-4\\1 \end{bmatrix}_2 (2^{N-4}+1) = 2^{2(N-4)} - 1$$
(22)

totally-isotropic PG(4, 2)s passing through a totally-isotropic PG(3, 2). And since the total number of PG(4, 2)s passing via a PG(3, 2) of PG(2N - 1, 2) is

$$2^{2N-4} - 1 \tag{23}$$

(as stemming from Eq. (5) for n = 2N, k = 5, l = 4 and q = 2), through a totally-isotropic PG(3, 2) there pass

$$2^{2N-4} - 1 - \left(2^{2(N-4)} - 1\right) = 15 \times 2^{2N-8}$$
(24)

PG(4, 2)s that feature totally-isotropic PG(3, 2)s apart from those that are totally isotropic. Hence, the number of those PG(4, 2)s of PG(2N - 1, 2) that contain totally-isotropic PG(3, 2)s – with the exclusion of totally isotropic ones – amounts to $(21) \times (24)/3$, where we also took into account (see Remark 5) that any such PG(4, 2) features just three totally-isotropic PG(3, 2)s. All in all, there are

$$\begin{bmatrix} 2N\\5 \end{bmatrix}_2 - \begin{bmatrix} N\\5 \end{bmatrix}_2 \prod_{i=1}^5 \left(2^{N+1-i}+1\right) - 15\begin{bmatrix} N\\4 \end{bmatrix}_2 2^{2N-8} \prod_{i=1}^4 \left(2^{N+1-i}+1\right)/3$$
(25)

PG(4, 2)s in the ambient PG(2N - 1, 2) that are not endowed with any totally-isotropic PG(3, 2)s, the number of quadratic doilies of W(2N - 1, 2) being just 16 times this number.

Employing further the fact that the three lines of a perp-set of a quadratic doily span a PG(3, 2), in particular that PG(3, 2) that contains just three totally-isotropic PG(2, 2)s, we find the compact formula

$$D_q(N) = \frac{48}{15} 4^{N-3} \Theta_3(N) \tag{26}$$

for the number of quadratic doilies in $\mathcal{W}(2N-1,2)$, where

$$\Theta_3(N) = \frac{7}{3} 2^{2N-6} \prod_{i=1}^3 \frac{2^{N-3+i} - 1}{2^i - 1} \prod_{i=1}^3 \left(2^{N+1-i} + 1 \right)$$
(27)

is the number of those PG(3,2)s of the ambient PG(2N-1,2) each of which features just three totally-isotropic PG(2,2)s.

N	$D_l(N)$	$D_q(N)$	D(N)
2	1	-	1
3	336	1 008	1 344
4	91 392	1370880	1462272
5	23744512	1495904256	1519648768
6	6100942848	1555740426240	1561841369088
7	1563272675328	1599227946860544	1600791219535872
8	400289425260544	1639185196441927680	1639585485867188224
9	102479956839235584	1678929132897196572672	1679031612854035808256

Table 1: First numbers D(N) (resp. $D_l(N)$, $D_q(N)$) of (resp. linear, quadratic) N-qubit doilies.

Comparing expressions (9) and (18), one gets

$$D_q(N) = (4^{N-2} - 1)D_l(N).$$
(28)

Consequently, the total number of doilies is

$$D(N) = 4^{N-2} D_l(N). (29)$$

For $2 \le N \le 9$ the numbers of N-qubit doilies are collected in Table 1. Since the quadratic doilies of $\mathcal{W}(2N-1,2)$ span a PG(4,2), it has no geometrical meaning to consider quadratic doilies for N = 2. It can nevertheless be noticed that Eq. (18) also holds for N = 2, and consistently gives $D_q(2) = 0$.

4 Generation of all *N*-qubit doilies

An N-qubit doily can be represented by an isomorphism f, sometimes called a (*doily*) labeling, mapping the points of W_2 to distinct points of W_N , preserving commutations and anticommutations,

and such that $f(a,b) = \pm f(a) \cdot f(b)$ for any two commuting points/observables a and b (the dot (.) denotes the matrix product). This section describes an algorithm for the enumeration of all N-qubit doilies, for any $N \ge 2$, by construction of one of their labelings.

Let us start with some definitions. An *N*-qubit ovoid is a 5-set of mutually anticommuting *N*qubit observables whose product is the identity \mathcal{I}_N . A triad is a 3-set of mutually anticommuting *N*-qubit observables. A center of a triad is a point commuting with the three points of the triad. A unicentric triad is a triad that has only one center. Let ε denote the empty word. The lexicographic order < on words is such that $\varepsilon < u$ for all non-empty word u, and a.u < b.v if and only if either a < b, or a = b and u < v, for any letters a and b and words u and v.

In order to avoid to consider several times objects that are similar but differently ordered, we define as follows a total order among letters and words, and then extend it to all tuples and sets of objects of the same nature, such as lines, sets of lines, etc. Pauli observables, encoded as words on the alphabet $\{I, X, Y, Z\}$, are totally ordered by the lexicographic order < induced by the order on letters, also denoted <, such that I < X < Z < Y. These orders are chosen so that their binary counterpart through the encoding $I \rightarrow 00, X \rightarrow 01, Z \rightarrow 10, Y \rightarrow 11$ is the lexicographic order on bit vectors (aka. bytes or binary words) induced by the order 0 < 1 on bits. This order < extends further to tuples (a_1, a_2, \ldots, a_n) of words, by considering them as words, by associating canonically to each set the tuple (a_1, a_2, \ldots, a_n) of its elements written in increasing order $(a_i < a_j \text{ when } i < j)$, and so on at any level of the hierarchy of objects of the same nature, such as a point-line geometries, seen as sets of lines, that are sets of points.



Figure 5: The 2-qubit doily W_2 , the ovoid O_2 (framed), the triad T_2 (framed and dashed), its center c_2 (circled and dashed) and the completion order (subscripted). The negative lines are doubled.

The algorithm relies on the following predefined elements, depicted in Fig. 5: the 2-qubit doily W_2 , the ovoid $O_2 \equiv \{IX, IZ, XY, ZY, YY\}$ in W_2 , the unicentric triad $T_2 \equiv \{IX, IZ, XY\}$ in O_2 , the center $c_2 \equiv XI$ of T_2 , and the sequence of lines

$$\begin{split} S &\equiv (XI, IX, XX), (XI, IZ, XZ), (XI, XY, IY), (ZY, XX, YZ), (ZY, XZ, YX), (ZY, IY, ZI), \\ & (YY, XX, ZZ), (YY, XZ, ZX), (YY, IY, YI). \end{split}$$

In the figure the third element of the tuples in this *completion order* is numbered from 1 to 9.

The algorithm itself is presented in Algorithm 1, where $f(a) \leftarrow b$ denotes the assignment of b as the image of a by f.

Algorithm 1 Doily generation algorithm.

1: for each ovoid $O = \{o_1, o_2, o_3, o_4, o_5\}$ in W_N , with $o_1 < o_2 < o_3 < o_4 < o_5$ do $f(IX) \leftarrow o_1 \parallel f(IZ) \leftarrow o_2 \parallel f(XY) \leftarrow o_3 \parallel f(ZY) \leftarrow o_4 \parallel f(YY) \leftarrow o_5$ 2: for each center c of $\{o_1, o_2, o_3\}$ in W_N that anticommutes with o_4 and o_5 do 3: $f(c_2) \leftarrow c$ 4: for each line (p, q, r) in the order of the sequence S do $f(r) \leftarrow |f(p).f(q)|$ end for 5:if O is not the smallest ovoid of f then discard f end if 6: \triangleright location for a potential treatment of f 7:end for 8: 9: end for

On Line 2 a doily labeling f is partially defined by the choice of images for the 5 points of the ovoid O_2 of W_2 . These images are the points of some ovoid $O = \{o_1, o_2, o_3, o_4, o_5\}$ of N-qubit observables. The points are assigned in increasing order so that to avoid duplicates. As these five assignments are independent, they can be performed in parallel.

Then (on Line 3) the algorithm looks for a point c that commutes with the first three points of O and that anticommutes with its last two points o_4 and o_5 . On Line 4 this point becomes the image by f of the center c_2 of the triad T_2 of O_2 .

The completion step on Line 5 computes one by one the images of all the other points of W_2 by f, in the order described by the sequence of lines S. At each iteration of this loop, for the line (p, q, r), the values f(p) and f(q) are known. By definition of a doily line, the image by f of the third point r is the product of the images f(p) and f(q) of the first two points, up to a possible minus sign, removed by the operation $|_|$ that denotes absolute value.

Knowing that each doily features 6 ovoids, the same doily is generated 6 times before Line 6, whose statement keeps only one of them, namely the doily d generated from the ovoid that is the smallest (according to the lexicographic order) among the 6 ovoids in d.

On Line 7 various treatments of the generated doilies f can be added, such as a storage, or the computation of classification criteria defined in Section 5.

4.1 Justification of the generation algorithm

First of all, the fact that doily labelings encode multi-qubit doilies is a direct consequence of the definition of a multi-qubit doily. Then, the properties of correctness and completeness for the doily enumeration algorithm mainly come from the following definition and proposition, whose proof is illustrated by Figure 6.

Definition 9 (Doily Root). Any pair (O, c) such that O is an ovoid of W_N and c is a point of W_N that commutes with exactly three points of O and anticommutes with the other two points is called a *(multi-qubit) doily root.*

Proposition 10. Any doily root (O, c) of W_N determines exactly one N-qubit doily.

Proof. Let $O = \{o_1, o_2, o_3, o_4, o_5\}$ and c be such that (O, c) is a multi-qubit doily root. The fact that c commutes with o_1 , o_2 and o_3 implies that there exist three multiplicative factors a_1 , a_2 and $a_3 \in \{-1, 1\}$ such that $\{c, o_i, a_i \ c. o_i\}$ are isotropic lines of W_N , for i = 1, 2, 3. This is depicted in Figure 6 by the 3 points $\pm c. o_i$. But because c anticommutes with o_4 and o_5 we also have



Figure 6: Expression of each observable according to the completion order followed by the generation algorithm, from the doily root ($\{o_1, o_2, o_3, o_4, o_5\}, c$). The thick lines are the lines used to compute these expressions.

that the observables $a_i \ c.o_i$ commute with o_4 and o_5 . Therefore there exist multiplicative factors $a_{ij} \in \{-1, 1\}$ such that $\{a_i \ c.o_i, o_j, a_{ij} \ (c.o_i).o_j\}$ are isotropic lines, for i = 1, 2, 3 and j = 4, 5. This is depicted in Figure 6 by the 6 points of the form $\pm (c.o_i).o_j$. These 9 points are computed by the completion step of the generation algorithm, in the order $a_1 \ c.o_1, a_2 \ c.o_2, a_3 \ c.o_3, a_{14} \ (c.o_1).o_4, a_{24} \ (c.o_2).o_4, a_{34} \ (c.o_3).o_4, a_{15} \ (c.o_1).o_5, a_{25} \ (c.o_2).o_5, a_{35} \ (c.o_3).o_5$. The 9 geometric lines thus identified are depicted by thick lines in Figure 6. Finally, it is easy to check that the six 3-sets represented by the thin lines in Figure 6 indeed are geometric lines. A noticeable property is that the product of the three points on each of these lines contains twice the center c and once each point of the ovoid O. By applying the known commutation and anticommutation relations between these points, it comes that the product of both centers annihilates. So, modulo a possible minus sign, it remains the product of all observables of the ovoid, known to equal identity. Therefore, the product of the three observables on each line equals $\pm \mathcal{I}_N$. Consequently, these 15 points and 15 lines form a doily, shown in Figure 6, so the algorithm is correct.

Each multi-qubit doily features at least one ovoid and the first loop explores all ovoids in W_N . So, each doily is found six times before Line 6, since each multi-qubit doily features exactly six ovoids. As the statement on this line always keeps one of them (the one that has been produced from the smallest of its ovoids), the algorithm is also complete.

4.2 Algorithmic complexity and implementation details

The enumeration algorithm explores all 4-tuples of observables likely to form an ovoid (the fifth point in the ovoid is computed as the product of the previous four), and then explores all observables to find c (on Line 3 of Algorithm 1). Therefore, the complexity of the algorithm is estimated to be $O(4^{5N})$, when the time unit is the duration to check whether two observables commute.

For efficiency reasons, we have implemented the algorithm in the C language, which allows for many optimizations. The total code is composed of about 2 300 lines and 50 functions, some of which implementing the classification process presented in Section 5. Some factors make the algorithm implementation more efficient than the former one presented in [7]: The new algorithm has a lower complexity; compared to the previously used language Magma [12], the low-level language C allows to perform fast operations on bit vector representations of the observables, using as few CPU instructions as necessary, and to split the workload into multiple threads.

The calculations were run on Linux Ubuntu, on a PC equipped with an Intel (R) Core(TM) i7-8665U 1.90 GHz and 15 GB RAM. The code was compiled with gcc 9.3.0 with optimization Ofast and is multi-threaded with OpenMP.

5 Multi-doily classification process and results

This section presents our classification criteria of N-qubit doilies and the classification results for N = 4 and N = 5.

5.1 Classification criteria

The classification parameters adopted are the same as in [7]. The classification of an N-qubit doily is based on: 1) its signature, i.e. the number of its observables containing a given number of I: $N-1, N-2, N-3, \ldots$ respectively named types $A, B, C, \ldots; 2$) the configuration of its negative lines, as described in Section 2; and 3) its linear or quadratic character.

To find the line configuration of a doily, the first discriminatory factor is the number of negative lines, since for each number of negative lines except 7 and 8, there is only one configuration possible. Then the property used to distinguish configurations 7A from 7B and 8A from 8B is to count the number of observables contained in at least one negative line, since this number is different between A and B.

We use the following property to check whether a doily is linear or quadratic. Given an N-qubit doily, we pick up in it a tricentric triad (here we take the image of $\{XY, ZY, YI\}$). If the product of the corresponding three observables is $\pm i\mathcal{I}_N$, then the doily is linear, otherwise it is quadratic. This is because any tricentric triad is a line in the ambient PG(3, 2) if a doily spans a PG(3, 2).

5.2 Database of numerical results

Using the program described in Section 4, we were able to classify all doilies for N = 3 (2016 ovoids), N = 4 (548352 ovoids) and N = 5 (142467072 ovoids). This classification is a treatment added on Line 7 of the algorithm presented in Algorithm 1, that determines the complete type of each generated doily, counts the number of doilies for each type, and registers it in a result table.

The sums of the numbers of linear and quadratic doilies found in each of the above-mentioned cases correspond exactly to those stemming from eqs. (9) and (18), respectively, summarized in

Table 1. The results of our classification are collected in Appendix A (three qubits), Appendix B (four qubits) and Appendix C (five qubits). The data for three qubits are in complete agreement with those of [7]; we found 11 different types of doilies of which five are linear and six quadratic. The 95 distinct types of four-qubit doilies split into 24 linear and 71 quadratic ones, whereas amongst 447 types of five-qubit doilies one finds 89 linear and 358 quadratic.

The structure of the classification table in each appendix is the same: the first column gives the type, the next N columns feature the numbers of observables of the corresponding types in a doily of the given type, the ν column shows the doily's character, and the remaining columns contain information about how many doilies of the given type are endowed with a particular number of negative lines (the blank space stands for zero here). The types are ordered in decreasing order of the number of observables containing no Is, in case of equality in decreasing order of the number of observables containing one I, and so on up to the number A of observables containing N - 1 Is. For instance, for 4 qubits, the type 1 contains the maximal number 12 of D-type observables, and the last type 95 contains only A- and B-type observables. For a given signature, the type of quadratic doilies precedes that of linear ones.

The result tables are stored in https://quantcert.github.io/.

The C code for classification runs in 0.3 s for 4 qubits with 1.4 MB of memory and 12 min with 1.8 MB of memory for 5 qubits. The memory usage is low because the doilies are not stored, all the measurements are performed on the fly.

5.3 Remarks about five-qubit doilies

Let us have a closer look at the five-qubit case. The $3^2 \times {5 \choose 2} = 90$ observables of type *B* and $3^4 \times {5 \choose 4} = 405$ observables of type *D* lie on an elliptic quadric $\mathcal{Q}^-_{(YYYYY)}(9,2)$ of $\mathcal{W}(9,2)$. This special quadric $\mathcal{Q}^-_{(YYYYY)}(9,2)$, like any non-degenerate quadric, is a geometric hyperplane of $\mathcal{W}(9,2)$. As a doily is also a subgeometry of $\mathcal{W}(9,2)$, it either lies fully in $\mathcal{Q}^-_{(YYYYY)}(9,2)$ (in which case $B \cup D = 15$, such a doily will be called special), or shares with $\mathcal{Q}^-_{(YYYYY)}(9,2)$ a set of points that form a geometric hyperplane, in particular an ovoid $(B \cup D = 5)$, a perp-set $(B \cup D = 7)$ and/or a grid $(B \cup D = 9)$ and being referred to as ovoidal, perplai and/or gridal, respectively.

From Appendix C one can infer a number of interesting properties. We first notice that signatures with $B \cup C$ being even or odd are endowed with even or odd numbers of negative lines, respectively.

We also observe that there are 12 different signatures with A = C = E = 0, i.e., signatures featuring solely special doilies.

Further, there are 17 particular signatures such that each features observables of every type and no two types have the same cardinality. Out of them, six are ovoidal (e.g., 2-1-3-4-5), seven perpial (e.g., 1-3-2-4-5) and four gridal (e.g., 2-4-3-5-1).

If all doilies of a particular signature have just five or just six negative lines, then each doily is ovoidal; if a signature features just seven negative lines, then all of its doilies are perpial.

Among 33 distinct signatures with four negative lines only, one finds 12 ovoidal, 11 perpial and 9 gridal ones; doilies of the remaining signature, viz. 0-8-0-7-0, are special.

Next, there are 15 different signatures whose doilies are endowed with 12 (i.e., the maximum number of) negative lines. Out of them, five are ovoidal, five perpial and four gridal; the doilies of the remaining signature, namely 0-0-0-15-0, are special. Similarly, there are 35 distinct signatures whose doilies contain 11 (i.e., the maximum odd number of) negative lines; out of them, 10 are

ovoidal, 11 perpial and 12 gridal, with the remaining two signatures, viz. 0-1-0-14-0 and 0-3-0-12-0, featuring solely special doilies.

5.4 Specific behavior of linear doilies

Finally, this section and the next one briefly mention some properties of linear doilies. Like the three- and four-qubit cases, a linear five-qubit doily can be either ovoidal or gridal and always contains an odd number of negative lines. Also, 75 types of linear doilies share their signatures with their quadratic siblings. However, there are 14 different signatures that are genuinely linear, of which eight cases are ovoidal.

From our results on three-, four- and five-qubit cases it follows that a linear doily (a) always features an odd number of negative lines, and (b) does not share a perp-set with the distinguished quadric.

We conjecture that Property (a) holds for any number of qubits $N \ge 2$, but we have not yet found of proof of it; we surmise that it has something to do with the fact that a linear doily is "squeezed" into a PG(3, 2), compared to a quadratic doily that enjoys more degrees of freedom being stretched out in a PG(4, 2). Property (b) can readily be proved to hold for any $N \ge 3$, as follows.

Proof. Let us consider a linear doily with one of its perp-set; on Figure 7a, this perp-set is illustrated in bold font. Any perp-set of any doily features four tricentric triads; in our perp-set these triads are: $\{1, 2, 3\}$, $\{3, 4, 5\}$, $\{1, 5, 6\}$ and $\{2, 6, 4\}$. Now, we know that any tricentric triad of a linear doily corresponds to a non-isotropic line in the ambient projective space, the four lines plus the three (totally isotropic) lines of the perp-set forming a Fano plane in this space, which is illustrated in Figure 7b. The assumption that our perp-set also lies on the distinguished quadric would mean that the whole plane would lie in the distinguished quadric and so would be totally isotropic, a contradiction.



Figure 7: Graphical arguments for the property that a linear doily cannot share a perp-set with the distinguished quadric.

5.5 A distinguished hexad of (linear) doilies

One knows that given an ovoid, there is a unique linear doily containing this ovoid. Now, take any quadratic doily. As each of its six ovoids defines a unique linear doily, we have a unique hexad of doilies tied to each quadratic doily. This holds for any $N \ge 3$. Figure 8 illustrates this property for N = 3. It features a quadratic doily in the middle, its six ovoids depicted explicitly as pentads of points located on bold gray lines and the corresponding six linear doilies; for better readability, the points of the corresponding ovoids are illustrated by double-circles.



Figure 8: A particular hexad of linear doilies in the three-qubit symplectic polar space.

6 Conclusion

There are a number of intriguing extensions and generalizations of the ideas and findings presented in this paper. We shall mention a few of them.

An interesting situation that will be worth addressing occurs in the case of N = 4. Given a PG(3,2) of the ambient PG(7,2) of $\mathcal{W}(7,2)$, its polar space is another PG(3,2). Hence, PG(3,2)s in PG(7,2) come in polar pairs. Taking into account the fact that a non-isotropic PG(3,2) features a

unique linear doily of W(7, 2), the above property means that also linear doilies of W(7, 2) occur in pairs. That is, picking up any linear four-qubit doily, there exists a unique linear doily such that each of its 15 observables commutes with each observable of the selected doily. This observation raises several interesting questions. For example, it would be interesting to ascertain which signatures are/can be paired, or which cardinalities of negative lines can occur in such self-polar pairs; we have already checked by hand a few examples where both doilies in a pair have the same signature and feature the same number of negative lines. There are (see Appendix B) altogether 24 different signatures featured by linear four-qubit doilies. We can then create a graph on 24 vertices such that its two edges are connected if there exists a pair of linear doilies exhibiting the corresponding signatures; we can even add a weight to an edge showing how many pairs of doilies feature this particular pair of signatures. This graph, as it follows from the examples checked, will also have edges joining a vertex to itself when the two paired signatures are identical. So, being an interesting graph of its own, it will also reveal some finer traits of the relation between individual linear doilies in W(7, 2)!

A particular case deserving closer attention is N = 6. Here, let us formally view any six-qubit observable as a 'syntheme' partitioned into three two-qubit observables ('duads'). Given a partition, we find a set of linear doilies such that any doily in the set features 15 particular observables such that when restricted to the same duad we get a two-qubit doily; that is, any such doily can formally be regarded as being composed of three two-qubit doilies. Moreover, each partition features a prominent doily having all the three duads identical. The next worth-exploring case in this respect is N = 9, as $\mathcal{W}(17, 2)$ hosts not only composites comprising three doilies having the same number of qubits (namely three), but also those whose compounds feature different numbers of qubits (namely four, three and two).

Another prospective, but much more challenging, task will be to count and classify all rank-three spaces, W(5, 2)s, living in a particular W_N , for $N \ge 4$. The case N = 4 was already briefly examined in [7]. To address higher rank cases, we plan to employ the strategy that is the direct and natural generalization of the ovoid-based algorithm for doilies described in this paper. Geometrically, an N-qubit ovoid is a set of five points lying on a certain elliptic quadric of a PG(3, 2) in the ambient PG(2N - 1, 2). Hence, its analogue will be a set of 27 N-qubit observables lying on an elliptic quadric of PG(5, 2) in the ambient PG(2N - 1, 2), and a triad of the ovoid will have its counterpart in a quadratic doily located on the quadric. A root of an N-qubit W(5, 2) will thus comprise an elliptic quadric and an off-quadric point such that its associated observable commutes with each of the 15 observables of a doily located in the quadric and anticommutes with the remaining 12 observables. It is obvious that this task will require a more elaborate generation algorithm and a more complex computer code to be successfully accomplished.

Acknowledgments

This project is supported by the EIPHI Graduate School (contract ANR-17-EURE-0002). This work was also supported by the Slovak VEGA Grant Agency, Project # 2/0004/20. We thank our friend Zsolt Szabó for the help in preparation of several figures.

References

- J.J. Sylvester. XLIV. Elementary researches in the analysis of combinatorial aggregation. The London, Edinburgh, and Dublin Philosophical Magazine and Journal of Science, 24(159):285– 296, 1844. https://doi.org/10.1080/14786444408644856.
- [2] L. Cremona. Teoremi stereometrici dal quali si deducono le proprietà dell' esagrammo di Pascal (Stereometric theorems from which the properties of Pascal's hexagram are deduced). Atti dell'Academia dei Lincei (Proceedings of the Accademia dei Lincei), coi tipi del salviucci edition, 1877.
- [3] H. W. Richmond. The figure formed from six points in space of four dimensions. *Mathematische Annalen*, 53:161–176, 1900.
- [4] S. E. Payne and J. A. Thas. *Finite Generalized Quadrangles*. European Mathematical Society, 2009.
- [5] J. W. P. Hirschfeld and J. A. Thas. General Galois Geometries. Clarendon Press, 1991.
- [6] P. J. Cameron. *Projective and Polar Spaces*. University of London, Queen Mary and Westfield College, 1992.
- [7] M. Saniga, H. de Boutray, F. Holweck, and A. Giorgetti. Taxonomy of Polar Subspaces of Multi-Qubit Symplectic Polar Spaces of Small Rank. *Mathematics*, 9(18):2272, 2021. http: //dx.doi.org/10.3390/math9182272.
- [8] F. Holweck. Geometric constructions over C and F₂ for Quantum Information. In Quantum Physics and Geometry, volume 25 of Lecture Notes of the Unione Matematica Italiana, pages 87–124. Springer Nature, 2019. https://arxiv.org/abs/1810.04258.
- [9] B. Polster. A Geometrical Picture Book. Springer New York, 1998.
- [10] H. de Boutray, F. Holweck, A. Giorgetti, P.-A. Masson, and M. Saniga. Contextuality degree of quadrics in multi-qubit symplectic polar spaces. 2022. https://arxiv.org/abs/2105.13798.
- [11] M. De Boeck, M. Rodgers, L. Storme, and A. Švob. Cameron-Liebler sets of generators in finite classical polar spaces. *Journal of Combinatorial Theory Series A*, 167(C):340–388, 2019. http://dx.doi.org/10.1016/j.jcta.2019.05.005.
- [12] W. Bosma, J. Cannon, and C. Playoust. The Magma Algebra System I: The User Language. Journal of Symbolic Computation, 24(3):235-265, 1997. http://dx.doi.org/10.1006/jsco. 1996.0125.

A Taxonomy of 3-qubit doilies

	Ob	serva	bles		$\begin{tabular}{ c c c c c c c c c c c c c c c c c c c$											
Type	A	В	C	ν	3	4	5	6	7A	7B	8A	8B	9	10	11	12
1	1	5	9	q				108								
2	0	7	8	q					81							
3	2	5	8	l			162									
4	3	5	7	q		324										
5	0	9	6	l	9				27							
6	2	7	6	q	216		162									
7	4	5	6	l	54											
8	2	9	4	l	81											
9	4	7	4	q	81											
10	0	15	0	q	36											
11	6	9	0	l	3											

B Taxonomy of 4-qubit doilies

	C	bser	vabl	es					Con	figuratio	on of neg	ative lin	es				
Type	A	B	C	D	ν	3	4	5	6	7A	7B	8A	8B	9	10	11	12
1	0	3	0	12	q	216				648				648			
2	0	4	0	11	q				3888			3888					
3	0	5	0	10	q	972		1944		4860	1944			1944			
4	1	0	5	9	q	648								648			
5	3	0	3	9	l	144											
6	0	6	0	9	q		1296		5184								
7	0	1	6	8	q	972				3888						972	
8	1	1	5	8	q				7776								
9	2	1	4	8	q	1944		1944									
10	2	1	4	8	l	972					972						
11	0	7	0	8	q			1944		972							
12	0	2	6	7	q				15552			11664	19440				
13	1	2	5	7	q	7776		13608			15552			1944			
14	1	2	5	7	l	3888					7776						
15	2	2	4	7	q		11664						3888				
16	3	2	3	7	q	1944		1944									
17	0	8	0	7	q		3888										
18	0	1	8	6	q	648		3888		1944	15552			11664			
19	1	1	7	6	q		19440		18144				11664				
20	0	3	6	6	q	7452		21384		30132	46 6 56			8424			
21	0	3	6	6	l	2592		1944		4860	11664			648		324	
22	2	1	6	6	q	3888		9720			1944						

	Ο	bser	vabl	es	1				Cont	figuratio	n of neg	ative lin	es				
Type	A	B	C	D	ν	3	4	5	6	7A	7B	8A	8B	9	10	11	12
23	1	3	5	6	q		46656		32 400				11664				
24	2	3	4	6	q	11016		11664			4860						
25	2	3	4	6	l			3888									
26	3	3	3	6	q		3888										
27	4	3	2	6	l	324											
28	0	9	0	6	q	2700		1944		324							
29	0	0	10	5	q				1296				3888				
30	1	0	9	5	q			5832						1944			
31	1	0	9	5	l	648					1944						
32	0	2	8	5	q		7776		23 328			15552	19440				
33	1	2	7	5	q	15552		46656			19440			3888			
34	0	4	6	5	q		11664		19 4 4 0				7776				
35	2	2	6	5	q		31104						3888				
36	1	4	5	5	q	7776		11664			3888						
37	1	4	5	5	l			7776									
38	2	4	4	5	q		7776										
39	0	1	10	4	q					1944	3888						
40	0	1	10	4	l			1944						1944			
41	1	1	9	4	q		7776		7776				3888				
42	0	3	8	4	q	6480		17496		23328	15552			3888			
43	2	1	8	4	q	1944					1944						
44	2	1	8	4	l			1944									
45	1	3	7	4	q		27216		7776				3888				
46	0	5	6	4	q	23328		23328		14580	13608						
47	0	5	6	4	l	972		7776		2916							
48	2	3	6	4	q	8424		15552									
49	1	5	5	4	q		31104		3888								
50	3	3	5	4	q		7776										
51	2	5	4	4	q	3888		3888									
52	2	5	4	4	l	1458											
53	4	5	2	4	q	486											
54	0	11	0	4	q	972											
55	0	2	10	3	q		3888		6480			3888	3888				
56	1	2	9	3	q	4536		3888									
57	1	2	9	3	l			7776									
58	0	4	8	3	q		15552		24 624			3888					
59	2	2	8	3	q		15552										
60	1	4	7	3	q	28512		27216			3888						
61	3	2	7	3	l	648											
62	0	6	6	3	q		19 440		3888								
63	2	4	6	3	q		19440										
64	1	6	5	3	q	7776		3888									

	Ο	bser	vabl	es]				Con	figuratio	n of neg	ative lir	nes				
Type	A	B	C	D	ν	3	4	5	6	7A	7B	8A	8B	9	10	11	12
65	1	6	5	3	l	3888											
66	3	4	5	3	q	3888											
67	0	3	10	2	q					3888							
68	0	3	10	2	l	648				1944							
69	1	3	9	2	q		7776		1296								
70	0	5	8	2	q	9720		11664		3888							
71	2	3	8	2	q	5832		3888									
72	2	3	8	2	l	1944											
73	1	5	7	2	q		7776										
74	0	7	6	2	q	8748		5832		972							
75	0	7	6	2	l	1944											
76	2	5	6	2	q	5832											
77	2	7	4	2	q	1944											
78	1	4	9	1	q	3888		3888									
79	1	4	9	1	l	1944											
80	0	6	8	1	q		3888										
81	1	6	7	1	q	3888											
82	3	4	7	1	q	1944											
83	0	5	10	0	q	2592											
84	1	5	9	0	q				432								
85	0	7	8	0	q	6480				324							
86	2	5	8	0	l			648									
87	3	5	7	0	q		1296										
88	0	9	6	0	q	4266											
89	0	9	6	0	l	36				108							
90	2	7	6	0	q	864		648									
91	4	5	6	0	l	216											
92	2	9	4	0	l	324											
93	4	7	4	0	q	324											
94	0	15	0	0	q	144											
95	6	9	0	0	l	6											

C Taxonomy of 5-qubit doilies

Observables					Configu	ration of n	egative li	nes				
Type $A B C D E\nu$, 3	4	5	6	7A	7B	8A	8B	9	10	11	12
1 0 0 1 5 9 q		-		Ŭ	,	58 320	011	02	19440	10		
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	r			12960		00010			10 110			
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		58 320		233 280				233 280		116 640		19 440
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		00020	174960	200 200	116640	466 560		200 200	262 440	110010	116640	10 110
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	000040		114000	12960	110.040	400 000			202 440		110.040	
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		116 640		421 200			116640	174960		116 640		
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		110.040	29160	421 200	58 320	29160	110.040	114,500		110.040		
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$			29100	9720	00 0 20	29100						
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$				19440			58 320					
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$			58 320	13440	238 140	247 860	00.020		145 800			
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$			08 320	291 600	236 140	247 800	233 280	233 280	143 800	58 320		
			58 320	291000		58 320	233 280	233 280		38 320		<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		116 640	38 320	291 600		38320	174960	233 280		116640		
		110 040	240.000	291 000	405 790	707 200	174 900	233 280	010 400	110 040	F0 200	<u> </u>
14 0 1 2 4 8 q		F0 200	349920	000.000	495 720	787 320		F0 200	816 480		58 320	
		58 320	E0 200	233 280	977.090	400.000		58320	145 900		49.740	
			58 320		277 020	422 820			145 800		43740	
17 2 1 0 4 8 <i>l</i>		154000	19440	010.400			201.000	500.000		110.040		
18 0 2 2 3 8 q		174960	00.000	816 480			291 600	583200		116640		<u> </u>
19 1 2 1 3 8 <i>l</i>			38880									
20 0 4 0 3 8 9			- /	58 320			58320					
21 0 3 2 2 8 q			349920			116 640			58 320		29 160	
22 0 5 0 2 8 q					53460				29 160		14580	
23 0 5 2 0 8 q			14580									
24 0 7 0 0 8 q					7290							
25 0 0 1 7 7 q					349 920	174 960			174 960			
26 0 1 1 6 7 q		174960		991 440			524880	758160		291600		58320
27 0 0 3 5 7 q	68 0 4 0		670680		641520	1399680			816 480		408240	
28 1 0 2 5 7 q		291600		583200				291600				
29 0 2 1 5 7 q			758160		874800	1458000			1166400		58320	
$30 \ 0 \ 1 \ 3 \ 4 \ 7 \ q$		1574640		3888000			1458000	3382560		1341360		174960
31 1 1 2 4 7 q	233 280		1049760			583200			233 280			
32 1 1 2 4 7 l			174960						58 320			
33 0 3 1 4 7 9	1	233280		583200			291600	524880				
34 2 1 1 4 7 q	1	213840						58320				
35 0 2 3 3 7 9	554040		1195560		845640	2682720			729000		58320	i
36 0 2 3 3 7 l			320760						145800			
37 1 2 2 3 7 9	r	388800		233280				233280				
38 0 4 1 3 7 9	1		58320		116640	58320						
39 2 2 1 3 7 q			116640			58320						
40 3 2 0 3 7 9		38 880										
41 0 3 3 2 7 9		758160		583 200				233280		58 320		
42 0 5 3 0 7 9	1	29 160										
43 0 1 0 8 6 9			58320		335 340	131 220			233 280		14580	
44 0 0 2 7 6 9		524880		1516320			1224720	2 507 760		758 160		77 760
45 1 0 1 7 6 9			466560			466 560		-	136 080			<u> </u>
46 0 2 0 7 6 9		174960		174960			233 280	291 600		58 3 20		
$47 \ 2 \ 0 \ 0 \ 7 \ 6 \ q$		58 320		19440								
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$			2012040		4753080	7 231 680			3625560		1341360	
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		874 800		349 920				641 520			,	
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		0.2000	174960		184 680	58 320			272 160		29160	
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$			1,1000		24300	14 580					4860	
	0100				-1000	11000	I		I		1000	

		Obs	erval	bles					Configu	ration of n	egative li	nes				
152 21 0 6 6 400 1283.040 <td>Type</td> <td></td> <td></td> <td></td> <td>3</td> <td>4</td> <td>5</td> <td>6</td> <td></td> <td></td> <td></td> <td></td> <td>9</td> <td>10</td> <td>11</td> <td>12</td>	Type				3	4	5	6					9	10	11	12
53 00 4 5 6 0 1691280 157420 137400 1283040 309360 55 10 3 5 6 7 2000 2012040 1953720 437400 77810 1 56 0 2 5 6 0 22506480 589320 209520 788120 788100 1 5 77810 1 1 5 77810 1 1 5 77810 1		2 1	0 6	6 q	64800			_						-		
14 10 3 5 6 7 900 2012 040 1953 720 1373 700 7760 56 0 2 25 6 0 2560 080 5800 320 2099 520 4782 240 758 160 758 160 57 2 2 5 6 7 210 5 6 7 80 7 80 7 80 1 4 6 972 000 3285 520 2.624 400 660 320 2.935 440 437 400 466 560 291 60 610 11 13 4 6 972 400 383 20 176 900 2.216 160 210 13 466 560 291 60 58 320 167 23 174 960 174 960 174 960 174 960 174 960 160 160 16 36 16 40 174 960 160 174 960 160 174 960 160 160 16 16 16 16 16 16 16 16 16 16 16 16	53					1691280		3965760			1574640	4490640		1283040		369 360
	54						2012040			1953720			437 400			
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	55	10	3 5	6l			349 920						77 760			
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	56	0 2	2 5	6 q	r	2566080		5890320			2099520	4782240		758160		
99 10 1 5 6 6 9 10 1	57					291 600		58320								
	58	12	1 5	6 q	272160		563760			233280			58 320			
	59	30	1 5	6 l	6480											
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	60	$0 \ 1$	4 4	6 q			3265920		2624400	6619320					437400	
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	61			6 l	87 480		583200		145800	437400			466560		29160	
	62	1 1	3 4	6 q	1	3849120		1769040				2216160				
65 21 2 4 61 29 90 58 320 66 0 5 0 4 6 q 28320 29 160 58 320 174 90 67 23 3 6 q 100 123 83 6 q 100 174 90 10									1078920	2653560			806 760		58320	
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $				-			495720									
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $																
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $									87480	29160			58 320			
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $							19440									
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $						2 4 4 9 4 4 0		2896560			466 560	2 216 160		174960		
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $							1 253 880						204 120			
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $						F00.000		F 4 4 2 2 5		116 640		000.000				<u> </u>
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $			2 3	6 q	1			544 320								
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $						758 160	00100					174960				
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $									145.000	220 740			80.000			<u> </u>
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $									145 800	320 760						
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $						F0 990	29 160	110040				F0 900	9720			
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $						58 320	000 000	116 640	00.1.00	110.040		58 320				
$ \begin{array}{ c c c c c c c c c c c c c c c c c c c$									29 160							<u> </u>
$ \begin{array}{ c c c c c c c c c c c c c c c c c c c$										29 160						
$ \begin{array}{ c c c c c c c c c c c c c c c c c c c$						20.160	29 160									
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $						29100										
$ \begin{array}{ c c c c c c c c c c c c c c c c c c c$						58 320										
$ \begin{array}{ c c c c c c c c c c c c c c c c c c c$				-		00 020										
$ \begin{array}{ c c c c c c c c c c c c c c c c c c c$				-			14 580									
$ \begin{array}{ c c c c c c c c c c c c c c c c c c c$							14000		2/30							
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $							116 640			379.080			388 800		58 320	
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$						58 320	110.040	194400	024000	010 000		58.320	000000		00020	
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $											1574640			583 200		174 960
$ \begin{array}{ c c c c c c c c c c c c c c c c c c c$						000 -000	233 280				10,1010	1000210		000 -00		111000
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $									3440880				6 006 960		1749600	
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $								2799360								
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $							1982880		1 399 680	2741040			1545480		204 120	
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $				-												
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $				5 q	1	6415200		13 996 800			6765120	14463360		3265920		349 920
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $							4257360									
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $								2157840		-		758160		116 640		
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	97					699840										
	98						4694760		4257360	9185400			6356880		1399680	
101 0 2 3 5 5 q 2624400 6298560 3265920 7989840 2449440 116640 102 2 0 3 5 5 q 87480 145800 103 1 2 2 5 5 q 1399680 816480 349920 104 0 4 1 5 5 q 291600 408240 320760 612360	99	00	5 5	5 l	87 480		355752		204 120	554040			437 400		87480	
102 2 0 3 5 5 q 87.480 145.800								2566080				3324240				
103 1 2 2 5 5 q 1399680 816480 349920 349920 104 0 4 1 5 5 q 291600 408240 320760 612360							6298560		3265920	7989840			2449440		116640	
104 0 4 1 5 5 q 291 600 408 240 320 760 612 360 612 360							145800									
						1399680		816 480				349920				
105 2 1 5 5 9 174 960 233 280 58 320				-					320760							
	105	2 2	1 5	5 q	174 960		233280			58320						

Observables				Configu	ration of n	legative li	ines				
Type $A B C D E \nu =3$	4	5	6	7A	7B	8A	8B	9	10	11	12
$106 \ 0 \ 1 \ 5 \ 4 \ 5 \ q$	7 290 000		9477000			2974320	8 164 800		1516320		58 320
107 1 1 4 4 5 q 1807920		3965760			2507760			524 880			
108 1 1 4 4 5 l 116640					233 280						
109 0 3 3 4 5 q	2916000		2274480			116 640	1 108 080				
110 2 1 3 4 5 q	1137240		58320				233 280				
$111 \ 1 \ 3 \ 2 \ 4 \ 5 \ q \ 262 \ 440$		320 760			58 320						
112 0 5 1 4 5 q	58 320		174960								
113 0 2 5 3 5 q 1399680		3819960		1 0 2 0 6 0 0	3674160			758 160		58320	
$114 \ 0 \ 2 \ 5 \ 3 \ 5 \ l \ 116 \ 640$				116 640	349 920						
115 1 2 4 3 5 q	2 332 800		933120				699840				
116 0 4 3 3 5 q 524880		816 480		58 320	349 920						
$117 \ 2 \ 2 \ 3 \ 3 \ 5 \ q \ 233 \ 280$		291600			58 320						
$118 \ 0 \ 3 \ 5 \ 2 \ 5 \ q$	758 160		583 200				174960				
$119 \ 1 \ 3 \ 4 \ 2 \ 5 \ q \ 58 \ 320$		116 640			58 320						
		58 320									
121 0 5 3 2 5 q	116 640										
122 2 3 3 2 5 q	58 320										
$123 \ 0 \ 4 \ 5 \ 1 \ 5 \ q \ 58 \ 320$		58320									
	1 224 720		2332800			1574640	3615840		1166400		349920
$125 \ 1 \ 0 \ 1 \ 9 \ 4 \ q \ 291 \ 600$		495 720			670 680			233 280			
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		174 960						58 320			<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		3645000		4 4 3 2 3 2 0	9127080			4403160		1210140	
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		466 560		306 180	714 420			466 560		43 740	
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	1 224 720	100 000	1 166 400	000100	111120		524 880	100 000		10 1 10	
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	1221120	116 640	1 100 100				021000				
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		110010			14 580						
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	5 598 720		11586240			5598720	14288400		4607280		933120
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		6 356 880	11 000 210		5 190 480			1516320			000120
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	4 315 680	0000000	6 706 800				4665600	1010020	291 600		
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	174 960		58 320			2 110 110	1000 000		201000		
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	111000	233 280	00010		116 640						<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		200 200			58 320						
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		15163200		10249740	26448120)		9856080		1501740	
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	7 290 000		4 121 280	10210110	20 110 120		3 324 240	0000000		1 001 1 10	
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		2595240	1121 200	1822500	2813940		0021210	787 320		14 580	
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		2000210		102060	466 560			101020		14 580	
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		1078920		102000	466 560					11000	
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	174 960	1010020	58 320		100 000						
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	111000		00020		14 580						
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	4782240		7678800			3557520	9506160		2041200		349920
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		3440880			2157840		0 000 100	524 880	- 0 11 - 00		010020
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		5 110 000			262 440			521000			<u> </u>
	11838960		14171760			2 099 520	7523280	<u> </u>	583 200		<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		4 082 400			2 041 200		. 020200	233 280	303 200		<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	1574640	1002 400	583 200		2011200		291 600	200 200			<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	816 480		000 200				174960				<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	010100	58 320					1.1000				<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		7771140		4 228 200	10585080			3207600		204 120	<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		116 640		204 120	699 840			58 320		29 160	\vdash
154 0 1 0 4 4 q 155 1 1 5 4 4 q	6 356 880	110010	2 361 960		000010		1866240	00020		-0 100	<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		4 4 3 2 3 2 0		1 851 660	4 1 26 1 40		- 000 - 10	437 400		14580	<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		408 240		- 001 000	145 800			10, 100		11000	\vdash
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		43740			110000	<u> </u>					\vdash
[조포포 [프] 포] 포] 포] 포] 인											
159 1 3 3 4 4 q	1 399 680	10 1 10	583 200				233 280				

Observables				Configu	ration of n	egative li	nes				
$TypeABCDE\nu = 3$	4	5	6	7A	7B	8A	8B	9	10	11	12
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		349 920	-	160 380	174 960		-	-			
$161 \ 0 \ 5 \ 2 \ 4 \ 4 \ l \ 7290$				21870							
162 2 3 2 4 4 q 233 280		160 380			58 320						
$163 \ 2 \ 5 \ 0 \ 4 \ 4 \ l \ 2430$											
164 0 2 6 3 4 q	3849120		2 332 800			349 920	1341360		58320		
$165 \ 1 \ 2 \ 5 \ 3 \ 4 \ q 1 \ 108 \ 080$)	1749600			729 000			87 480			
		116 640									
167 0 4 4 3 4 q	1924560		466 560				174960				
$168 \ 2 \ 2 \ 4 \ 3 \ 4 \ q$	524 880						58 320				
169 1 4 3 3 4 q 19440		58 320									
$170 \ 0 \ 6 \ 2 \ 3 \ 4 \ q$	116 640										
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		466 560		87 480	174 960						
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		58 320									
173 1 3 5 2 4 q	379 080		116 640								<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	0.0000	145 800	110010	160 380	131 220						<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		58 320									<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		00010								1	<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	58 320										<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	00020	58 320								1	<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		29160									
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		20100									
181 0 7 4 0 4 q 7290											<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$											
182 0 9 2 0 4 <i>q</i> 1230 183 0 0 3 9 3 <i>q</i> 907 200		3 0 3 2 6 4 0		1 720 440	5 452 920			3 936 600		1720440	<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		3 0 3 2 0 4 0		194 400	612 360			5 950 000		194400	<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	2 216 160		1 185 840	194400	012 300		1 982 880			194 400	<u> </u>
185 1 0 2 9 3 q 186 0 1 3 8 3 q	6 881 760		11537360 11547360				12013920		1924560		174960
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	0 881 700	2 332 800	11 547 500		1516320	4 102 240	12 013 920	233 280	1 924 000		174 900
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		2 332 800			233 280			233 280			<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	524 880				200 200		174960				<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		7115040		4 215 620	14929920			7153920		1982880	
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	4782240	7 110 040	3 304 800	4313080	14 929 920		2 799 360	1 100 920		1 982 880	<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		7727400	3 304 800	4 082 400	9 593 640			2536920		145 800	<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		174 960		233 280	1 108 080			$\frac{2330920}{58320}$		140 800	
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		174900 233 280		233 280	1 108 080			38 3 20			<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	2157840	200/200	758 160				641 520				<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	2107 840	58 320	758100				041 520				<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	13471920		20 256 480			6 500 160	15396480		1458000		174960
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		8019000			5015520	0.030.100	10 090 480	699840	1 400 000		114 900
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	4918320	0.019.000	4 335 120		0010020	174.060	1749600	033 040	58 320		<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	4 918 320 1 166 400		+ 555 120			114900	233 280		00 0 20		<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	1 100 400	524 880			116 640		200200				<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		116640			110.040						<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$				2507760	6 006 960			0 005 000		174.060	<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		3586680 174960		2507760	291 600			2225880 136080		174960	<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	1 399 680	114 900	1 283 040	110.040	291 000		583 200	100.000			<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		10847520		3849120	9622800		000 200	1458000			<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	5 307 120	10047020	1 632 960	0.049120	044 000		1 224 720	1 400 000			<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		1778760	1 052 900	116 640	524880		1 224 1 20				<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		349 920		110.040	024 000						<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$					174960						<u> </u>
	291 600	816 480			174900						┝───┤
211 1 4 2 5 3 q											┝───┤
212 3 2 2 5 3 q	116 640		1								
213 0 1 7 4 3 q	6 123 600		6862320			1 007 000	3615840		116640		

Configuration of regaring intersecture Configuration of regaring intersecture Viscal is 1 is 1 is 4 is 4	Observables					Configu	ration of n	egative li	ines				
2141 11 6 4 3 2 216 0 11 6 4 1 11 6 11 6 11 6 11 6 11 6 11 6 11		3	4	5	6	~		~		9	10	11	12
215 1 6 4 1 6 7 5 75 500 75			-		Ŭ			011	02		10		
1216 03 5 4 3 q 5715 300 3265 920 174 900 150 800 58320 0 0 0 218 13 4 3 q 854 30 1399 60 0 349 920 0 583 43 10 0 <td></td>													
217 21 1 5 4 3 q 758 160			5715360		3 265 920			174960	1 108 080				
118 13 4 3 6 84920 34920 0 0 0 0 0 120 0 3 4 3 0													
191 11 4 3 l		884520		1399680			349 920						
211 21 3 4 3 q 466 500 m													
211 21 3 4 3 q 466 500 m	220 0 5 3 4 3 q		1166400		174960								
1233 3 2 4 3 7 90 1 <td></td> <td></td> <td>466560</td> <td></td>			466560										
1221 12 7 3 3 1973160 2507760 1137240 1574640 174960 1 225 12 6 3 1 29160 233280 87480 1 <t< td=""><td>222 1 5 2 4 3 1</td><td>29160</td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td></t<>	222 1 5 2 4 3 1	29160											
225 0 2 7 3 1 291 0 1 1 1 1 0 1 1 0 1 1 0 1 <td>223 3 3 2 4 3 q</td> <td>29160</td> <td></td>	223 3 3 2 4 3 q	29160											
225 0 2 7 3 1 291 0 1 1 1 1 0 1 1 0 1 1 0 1 <td>224 0 2 7 3 3 q</td> <td>1973160</td> <td></td> <td>2507760</td> <td></td> <td>1 137 240</td> <td>1574640</td> <td></td> <td></td> <td>174960</td> <td></td> <td></td> <td></td>	224 0 2 7 3 3 q	1973160		2507760		1 137 240	1574640			174960			
1272 0 4 5 3 1/2 82592 1 4 5 3 1/1 116640 174960 0 <td></td> <td></td> <td></td> <td>233280</td> <td></td> <td>87 480</td> <td></td> <td></td> <td></td> <td></td> <td></td> <td></td> <td></td>				233280		87 480							
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $			1749600		116 640				233 280				
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	227 0 4 5 3 3 q	835 920		1283040		116 640	174 960						
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $				204120									
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $			699840		116 640								
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	230 0 6 3 3 3 q	116640		58320									
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $			524880		213840			58320					
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $		291600		291600			58 320						
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $		29160											
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $			524880		58320								
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	235 2 3 5 2 3 q		145800										
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	236 1 5 4 2 3 q	58320		29160									
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	237 3 3 4 2 3 q	29160											
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	238 0 4 7 1 3 q	58320											
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $		58320											
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	240 0 0 4 9 2 q		1283040		3168720			1166400	3615840		1049760		116640
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	241 1 0 3 9 2 q	602640		1487160						456840			
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $													
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $						4126140				3557520		262440	
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $		233280		262440		320760	933 120			204120		29160	
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	245 1 1 3 8 2 q		4082400		2 332 800				1516320				
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	246 2 1 2 8 2 q	233280					174960						
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$				58320									
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$			1807920		3868560			1866240	3790800		408240		
$\begin{array}{c c c c c c c c c c c c c c c c c c c $		699840		2216160			991440			291600			
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$					7523280			1458000	3 090 960		116640		
$\begin{array}{c c c c c c c c c c c c c c c c c c c $			116640										
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		1312200					787320			87480			
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	$\begin{array}{c c c c c c c c c c c c c c c c c c c $			174960									
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$			524880										
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$				$892296\overline{0}$		6765120	11401560			2303640		$116\overline{640}$	
$\begin{array}{c c c c c c c c c c c c c c c c c c c $			6648480		1749600				1166400				
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$							2478600			262 440			
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$						102 060							
261 2 3 2 6 2 q 58 320 58 323 323 323 320 50 52 <td></td> <td>481140</td> <td></td> <td>437400</td> <td></td> <td></td> <td>174960</td> <td></td> <td></td> <td></td> <td></td> <td></td> <td></td>		481140		437400			174960						
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $			$122472\overline{0}$		77 760								
$ \begin{array}{c c c c c c c c c c c c c c c c c c c $				58320									
264 1 0 7 5 2 q 379 080 699 840 233 280		$14\overline{580}$											
265 1 0 7 5 2 1 116			874 800		1283040			641520	699 840				
266 0 2 6 5 2 q 9097920 6531840 816480 1924560	264 1 0 7 5 2 q	379080					233 280						
	265 1 0 7 5 2 <i>l</i>			116640									
267 1 2 5 5 2 1 1 5 2 880 1 <td></td> <td></td> <td>9097920</td> <td></td> <td>6 531 840</td> <td></td> <td></td> <td>816 480</td> <td>1924560</td> <td></td> <td></td> <td></td> <td></td>			9097920		6 531 840			816 480	1924560				
	267 1 2 5 5 2 a	2216160		2857680			524880						

Observables					Configu	ation of n	egative li	nes				
Type $A B C D E \nu$	3	4	5	6	7A	7B	8A	8B	9	10	11	12
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		2 624 400	0	583 200	111		011	58 320	5	10	- 11	12
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		641520		000 200				00020				i
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		041020	116 640									<u> </u>
			110 040									
			0.440.440		1 1 97 9 40	1 510 000			000 000			
272 0 1 8 4 2 q			2 4 4 9 4 4 0		1 137 240	1516320			233 280			
273 0 1 8 4 2 <i>l</i>			291600		87480							
$274 \ 1 \ 1 \ 7 \ 4 \ 2 \ q$		2187000		233280				58320				
275 0 3 6 4 2 q			3863700		1166400	1370520			29160			
$276 \ 2 \ 1 \ 6 \ 4 \ 2 \ q$	116640		116640									
$277 \ 2 \ 1 \ 6 \ 4 \ 2 \ l$	14580											
278 1 3 5 4 2 q		2099520		116640								
279 0 5 4 4 2 q	575910		466560		123930							
280 0 5 4 4 2 l	87 480											
281 2 3 4 4 2 q	204 1 20		58320									
282 2 5 2 4 2 q												
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		2099520		1049760			116640					
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$			758160	0.20.00		116 640						
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$												
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		1 108 080		194 400								
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		349 920		101100								
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		040 020	116 640									il
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$			262 440		29160							
			202 440		29100							
290 0 3 8 2 2 <i>l</i>		174.000										<u> </u>
291 1 3 7 2 2 q		174960	1.45.000		20.1.00							
292 0 5 6 2 2 q			145800		29160							<u> </u>
293 2 3 6 2 2 q												
294 0 7 4 2 2 q												
295 2 5 4 2 2 q												
296 0 0 5 9 1 q			145800		291600	991440			145800		58320	
297 0 0 5 9 1 l			116640						116640			
298 1 0 4 9 1 q		699840		291600				291600				
$299 \ 2 \ 0 \ 3 \ 9 \ 1 \ q$			116640									
$300 \ 0 \ 1 \ 5 \ 8 \ 1 \ q$		2216160		2857680			1283040	2274480				
301 1 1 4 8 1 q	933120		1691280			758160			116640			
302 1 1 4 8 1 l			233280									
303 2 1 3 8 1 q		583200						58320				
304 0 0 7 7 1 q	77 760		466560		291600	349920			174960			
305 1 0 6 7 1 q		466560		272160				116640				
306 0 2 5 7 1 q			4694760		1895400	2303640			291600			
307 0 2 5 7 1 1			466560		174960							
308 1 2 4 7 1 q		2566080		291 600				291600				
$309 \ 2 \ 2 \ 3 \ 7 \ 1 \ q$			233 280									
$310 \ 3 \ 2 \ 2 \ 7 \ 1 \ q$		58 320										<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		2449440		2916000			699840	583 200				
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		_ 110 110	1 341 360	2010000		291 600	300010	300 200				
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		4607280	- 0 11 000	1691280		_01000	233 280	233 280				
314 2 1 5 6 1 q		466 560		1001200			200 200	200 200				
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		100 000	991440			116 640						
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$			331 440			110.040						├───┤
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		233 280										<u> </u>
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		200 200	20.160		174960	59 200			20.160			<u> </u>
			29160			58320			29 160			
319 0 0 9 5 1 l			9 790 400		29160	022 100						
320 0 2 7 5 1 q		0.000.000	3732480	000.000	1 399 680	933 120						
321 1 2 6 5 1 q		2332800		233280								

Observables				Configu	ration of n	egative li	nes				
$TypeABCDE\nu 3$	4	5	6	7A	7B	8A	8B	9	10	11	12
322 0 4 5 5 1 q 1574640		1 341 360	-	58 320	116 640	-	-	-			
323 0 4 5 5 1 l 174960											
324 2 2 5 5 1 q 291600		116640									
325 1 4 4 5 1 q	583200										
326 2 4 3 5 1 q 58 320											
327 0 1 9 4 1 q	991440		622 080			58 320	58320				
328 1 1 8 4 1 9 349 920		291600									
329 1 1 8 4 1 1 58 320											
330 0 3 7 4 1 q	2507760		408 240								
331 1 3 6 4 1 9 874 800		262440									
332 0 5 5 4 1 q	349920										
333 1 5 4 4 1 q 145 800											
334 3 3 4 4 1 q 29160											
335 0 2 9 3 1 q 913680		787320		204 1 20	58 320						
336 0 2 9 3 1 <i>l</i> 58 320											
337 1 2 8 3 1 q	233280										
338 0 4 7 3 1 q 816 480		408240		58 320							
339 2 2 7 3 1 q 116 640											
340 0 6 5 3 1 q 58 320											
341 0 3 9 2 1 q	87480										
342 1 3 8 2 1 q 116 640											
343 1 5 6 2 1 q 58 320											
344 0 0 0 15 0 q	311040		155520			136080	641520		186624		103680
345 0 1 0 14 0 q 534 600		1603800		831 060	2741040			1759320		554040	
346 0 2 0 13 0 q	3207600		4432320			1399680	3965760		349920		
347 0 3 0 12 0 q 2 227 500		3936600		1383480	4169880			926640		29160	
348 0 4 0 11 0 q	3790800		2099520			136080	874800				
349 0 5 0 10 0 q 1 526 040		1648512		330 480	592920			9720			
350 0 6 9 0 q	58320		77760			58320	58320				
$351 \ 1 \ 0 \ 5 \ 9 \ 0 \ q \ 12960$								3240			
352 1 0 5 9 0 l		58320									
353 2 0 4 9 0 q	58320										
354 3 0 3 9 0 <i>l</i> 3960											
355 0 6 0 9 0 q	939600		142560								
356 0 1 6 8 0 q 165240		174960		311 040	262440			58320		4860	
357 0 1 6 8 0 <i>l</i> 14580				43740							
358 1 1 5 8 0 q	466560		97200								
359 2 1 4 8 0 q 38 880		9720									
360 2 1 4 8 0 <i>l</i> 19440					4860						
361 0 7 0 8 0 q 340 200		126360		19 440							
362 0 0 8 7 0 q	58320		97 200			58320	58320				
363 1 0 7 7 0 q 58320		58320									
364 0 2 6 7 0 q	816480		777600			174960	97200				
365 1 2 5 7 0 q 505 440		476280			136 080			9720			
366 1 2 5 7 0 <i>l</i> 77760					38 880						
367 2 2 4 7 0 q	174960						19440				
368 3 2 3 7 0 <i>q</i> 38880		9720									
369 0 8 0 7 0 q	19440										
370 0 1 8 6 0 q 304560		281880		213 840	106 920			58320			
371 1 1 7 6 0 q	213840		90720				58320				
372 0 3 6 6 0 q 1 373 760		952560		340 200	291 600			42120			
373 0 3 6 6 0 <i>l</i> 81000		9720		24 300	58320			3240		1620	
374 2 1 6 6 0 q 77760		77760			9720						
375 1 3 5 6 0 q	349920		162000				58320				

Observables	Observables Configuration of negative lines										
$TypeABCDE\nu = 3$	4	5	6	7A	7B	8A	8B	9	10	11	12
376 2 3 4 6 0 q 842		58 320	Ŭ		24300	011	012		10		
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	10	19440			-1000						
378 3 3 3 6 0 q	19 440										
379 4 3 2 6 0 1 16											
380 0 9 0 6 0 q 55 0	080	9720		1620							
381 0 0 10 5 0 q			6480				19440				
382 1 0 9 5 0 q		29160						9720			
383 1 0 9 5 0 1 129	960				9720						
384 0 2 8 5 0 q	855 360		233 280			77 760	97200				
385 1 2 7 5 0 q 427	680	408 240			97200			19440			
386 0 4 6 5 0 q	349 920		97 200				38 880				
387 2 2 6 5 0 q	155 520						19440				
388 1 4 5 5 0 q 184		58320			19440						
389 1 4 5 5 0 1		38 880									
390 2 4 4 5 0 q	38 880										
391 0 1 10 4 0 q 252	720	87 480		9720	19440						
392 0 1 10 4 0 <i>l</i>		9720						9720			
393 1 1 9 4 0 q	38 880		38 880				19440				
394 0 3 8 4 0 91577	880	349 920		160 380	77760			19440			
395 2 1 8 4 0 q 97					9720						
396 2 1 8 4 0 <i>l</i>		9720									
397 1 3 7 4 0 q	136 080		38 880				19440				
398 0 5 6 4 0 q 750	870	116640		72900	68040						
399 0 5 6 4 0 <i>l</i> 48	60	38 880		14580							
400 2 3 6 4 0 q 110	160	77 760									
401 1 5 5 4 0 q	155 520		19440								
402 3 3 5 4 0 q	38 880										
403 2 5 4 4 0 9 194	40	19440									
404 2 5 4 4 0 1 72	90										
405 4 5 2 4 0 9 24	30										
406 0 11 0 4 0 q 48	60										
407 0 2 10 3 0 q	136 080		32400			19440	19440				
408 1 2 9 3 0 9 518	340	19440									
409 1 2 9 3 0 1		38880									
410 0 4 8 3 0 q	77 760		123 120			19440					
411 2 2 8 3 0 q	77 760										
412 1 4 7 3 0 9 142	560	136080			19440						
413 3 2 7 3 0 1 32	40										
414 0 6 6 3 0 q	97 200		19440								
415 2 4 6 3 0 q	97 200										
416 1 6 5 3 0 q 388		19440									
417 1 6 5 3 0 1 194											
418 3 4 5 3 0 q 194	140										
419 0 3 10 2 0 q 383	940			19440							
420 0 3 10 2 0 1 32				9720							
421 1 3 9 2 0 q	38 880		6480								
422 0 5 8 2 0 q 332		58320		19440							
423 2 3 8 2 0 q 29 1		19440									
424 2 3 8 2 0 1 97											
425 1 5 7 2 0 q	38 880										
426 0 7 6 2 0 q 437		29160		4860							
427 0 7 6 2 0 1 97											
428 2 5 6 2 0 q 29 1											
429 2 7 4 2 0 q 97	20										

Observables	Configuration of negative lines											
Type $A B C D E \nu$	· 3	4	5	6	7A	7B	8A	8B	9	10	11	12
430 1 4 9 1 0 9	19440		19440									
431 1 4 9 1 0 1	9720											
432 0 6 8 1 0 9	!	19440										
433 1 6 7 1 0 q	19 440											
434 3 4 7 1 0 9	9720											
435 0 5 10 0 0 q												
436 1 5 9 0 0 q	!			1080								
437 0 7 8 0 0 9	32 400				810							
438 2 5 8 0 0 1			1620									
439 3 5 7 0 0 q	!	3240										
440 0 9 6 0 0 9	21 330											
441 0 9 6 0 0 l	90				270							
442 2 7 6 0 0 q	2160		1620									
443 4 5 6 0 0 l	540											
444 2 9 4 0 0 1	810											
445 4 7 4 0 0 q	810											
446 0 15 0 0 0 q	360											
447 6 9 0 0 0 1	10											